Problem 1.

Let N be your University of Calgary ID number.

(a) Use the Euclidean Algorithm to compute gcd(N, 271) and use this to find integers x and y so that gcd(N, 271) = Nx + 271y.

Solution. (Answers will differ for different N.) Let's say that N = 12345678. Then

$$12345678 = 45556 \cdot 271 + 2$$

$$271 = 135 \cdot 2 + 1$$

$$2 = 2 \cdot 1 + 0.$$

This means that

$$gcd(12345678, 271) = gcd(271, 2)$$
$$= gcd(2, 1)$$
$$= gcd(1, 0) = 1.$$

Then

$$1 = 271 - 135 \cdot 2$$

= 271 - 135 \cdot (12345678 - 45556 \cdot 271)
= 271 + (135 \cdot 45556) \cdot 271 - 135 \cdot 12345678
= 271 + (6150060) \cdot 271 - 135 \cdot 12345678
= (-135) \cdot 12345678 + 6150061 \cdot 271,

so we can set x = -135 and y = 6150061 so that gcd(N, 271) = Nx + 271y. Or, using the "table method":

		x	y
R_1	12345678	1	0
R_2	271	0	1
$R_3 = R_1 - 45556 \cdot R_2$	2	1	-45556
$R_4 = R_2 - 135 \cdot R_3$	1	-135	$1 - 135 \cdot (-45556)$

and $1 - 135 \cdot (-45556) = 6150061$.

(b) Suppose that M is an integer such that gcd(M, 271) = gcd(M, 2016). Find gcd(M, 271). Explain how you get the answer.

Solution. gcd(M, 271) = 1.

Explanation. We note that 271 is prime, so its positive divisors are only 1 and 271, thus gcd(M, 271) can only be 1 or 271. Next we note that $2016 = 2^5 \cdot 3^2 \cdot 7$ and all of the divisors of 2016 are

 $\{ d \in \mathbb{Z}^+ \mid d \text{ divides } 2016 \} = \{ 1, 2, 3, 4, 6, 7, 8, 9, 12, 14, 16, 18, 21, 24, 28, 32, 36, 42, 48, 56, 63, \\72, 84, 96, 112, 126, 144, 168, 224, 252, 288, 336, 504, 672, 1008, 2016 \}.$

The only divisor that 271 and 2016 have in common is 1. Thus, the only way to have gcd(M, 271) = gcd(M, 2016) is when gcd(M, 271) = gcd(M, 2016) = 1.

Since gcd(K, 2016) must be a divisor of 2016 that satisfies 271 > gcd(K, 2016) > 250, the only divisor of 2016 that satisfies 271 > d > 250 is d = 252.

(c) Suppose that K is an integer between 800,000 and 900,000 so that gcd(K, 271) > gcd(K, 2016) > 250. Find all possible values of K. Explain how you get the answers. Solution. K = 887,796.

Explanation. From the condition that gcd(K, 271) > gcd(K, 2016) > 250 and from part (b), we conclude that gcd(K, 271) = 271 and gcd(K, 2016) = 252. Thus K must be a multiple of both 271 and 252, so K must be a multiple of $271 \cdot 252$. The only multiples of $271 \cdot 252$ that are between 800,000 and 900,000 are

 $12 \cdot 271 \cdot 252 = 819,504$ and $13 \cdot 271 \cdot 252 = 887,796.$

However, $gcd(12 \cdot 271 \cdot 252, 2016) = 1008$. So K cannot be $12 \cdot 271 \cdot 252$. Hence K must be equal to $13 \cdot 271 \cdot 252$.

Problem 2.

Consider the sequence of Fibonacci numbers f_1, f_2, f_3, \ldots which are defined as follows: $f_1 = 1, f_2 = 1$, and

$$f_n = f_{n-1} + f_{n-2}$$

for all integers $n \geq 3$.

(a) Prove that for all integers $n \ge 3$, $gcd(f_n, f_{n-1}) = gcd(f_{n-1}, f_{n-2})$. (You may use Lemma 4.8.2. No induction is needed.)

Proof. Let $n \geq 3$ be an arbitrary integer. Let $a = f_n$, $b = f_{n-1}$, $r = f_{n-2}$, and q = 1. Then

$$a = qb + r$$

since $f_n = f_{n-1} + f_{n-2}$ by definition of the sequence. From Lemma 4.8.2, we see that gcd(a, b) = gcd(b, r). Thus $gcd(f_n, f_{n-1}) = gcd(f_{n-1}, f_{n-2})$.

(b) Prove by weak induction that $gcd(f_n, f_{n-1}) = 1$ for all integers $n \ge 2$. (Use part (a).)

Proof. We prove this by induction.

Base case (n = 2): Note that $gcd(f_2, f_1) = gcd(1, 1) = 1$. Inductive step: Let $k \ge 2$ be an integer. Suppose that

$$gcd(f_k, f_{k-1}) = 1. \tag{IH}$$

(We want to show that $gcd(f_{k+1}, f_k) = 1$.) Now $k+1 \ge 2$, so using part (a) we see that

$$gcd(f_{k+1}, f_k) = gcd(f_k, f_{k-1})$$
 from part (a)
= 1 by IH,

which is what we wanted to show.

Therefore, by induction, $gcd(f_n, f_{n-1}) = 1$ for all integers $n \ge 2$

(c) Prove by weak induction that
$$\sum_{i=1}^{n} (f_i)^2 = f_{n+1}f_n$$
 for all integers $n \ge 1$.

Proof. We prove this by induction.

Base case (n = 1): Note that $\sum_{i=1}^{1} (f_i)^2 = (f_1)^2 = 1$ and $f_2 f_1 = 1 \cdot 1 = 1$. Hence $\sum_{i=1}^{1} (f_i)^2 = f_{1+1} f_1$.

Inductive step: Let $k \ge 1$ be an integer. Suppose that

$$\sum_{i=1}^{k} (f_i)^2 = f_{k+1} f_k.$$
 (IH)

(We want to show that $\sum_{i=1}^{k+1} (f_i)^2 = f_{k+2}f_{k+1}$.) Now

$$\sum_{i=1}^{k+1} (f_i)^2 = \sum_{i=1}^k (f_i)^2 + (f_{k+1})^2$$

= $f_{k+1}f_k + (f_{k+1})^2$ by IH
= $f_{k+1}(f_k + f_{k+1})$
= $f_{k+1}(f_{k+2})$ since $f_{k+2} = f_{k+1} + f_k$
= $f_{k+2}f_{k+1}$

which is what we wanted to show.

Therefore, by induction,
$$\sum_{i=1}^{n} (f_i)^2 = f_{n+1} f_n$$
 for all integers $n \ge 1$.

Problem 3.

The sequence s_0, s_1, s_2, \ldots , is defined by $s_0 = 1$, and for all integers n > 0,

$$s_n = s_{\lfloor \frac{n}{2} \rfloor} + s_{\lfloor \frac{2n}{3} \rfloor} + n.$$

The sequence t_0, t_1, t_2, \ldots , is defined by $t_0 = 2, t_1 = 3$, and for all integers $n > 0, t_n = 3t_{n-1} - 2t_{n-2}$.

(a) Find $s_1, s_2, s_3, s_4, s_5, s_6, s_7$, and s_8 . Guess the smallest integer a so that $s_n > 4n$ for all integers $n \ge a$. Solution: We see that

$s_1 = s_{\lfloor \frac{1}{2} \rfloor} + s_{\lfloor \frac{2}{3} \rfloor} + 1 = s_0 + s_0 + 1 = 1 + 1 + 1$	=3
$s_2 = s_{\lfloor \frac{2}{2} \rfloor} + s_{\lfloor \frac{4}{3} \rfloor} + 2 = s_1 + s_1 + 2 = 3 + 3 + 2$	=8
$s_3 = s_{\lfloor \frac{3}{2} \rfloor} + s_{\lfloor \frac{6}{3} \rfloor} + 3 = s_1 + s_2 + 3 = 3 + 8 + 3$	=14
$s_4 = 20$	
$s_5 = 27$	
$s_6 = 40$	
$s_7 = 41$	
$s_8 = 55,$	

so we can guess that $s_n > 4n$ for all $n \ge 3$.

(b) Prove by strong induction that $s_n > 4n$ for all integers $n \ge a$, where a is the integer you chose in part (a).

Proof. We prove this by induction.

Base cases:

(n = 3): Note that $s_3 = 14$ and $4 \cdot 3 = 12$. Thus $s_3 > 4 \cdot 3$ since 14 > 12. (n = 4): Note that $s_4 = 20$ and $4 \cdot 4 = 16$. Thus $s_3 > 4 \cdot 4$ since 20 > 16.

(n = 5): Note that $s_5 = 27$ and $4 \cdot 5 = 20$. Thus $s_4 > 4 \cdot 5$ since 27 > 20.

Inductive step: Let $k \ge 5$ be an integer. Suppose that

$$s_i > 4i$$
 for all integers $i \quad 3 \le i \le k$. (IH)

(We want to show that $s_{k+1} > 4(k+1)$.) First note that $\lfloor x \rfloor + 1 > x$ is true for all $x \in \mathbb{R}$ by definition of the floor. Thus

$$\lfloor x \rfloor > x - 1$$
 for all $x \in \mathbb{R}$. (*)

Furthermore, note that

$$3 \le \left\lfloor \frac{k+1}{2} \right\rfloor \le k$$
 and $3 \le \left\lfloor \frac{2k+2}{3} \right\rfloor \le k$ are both true since $k \ge 5$, (**)

so we may use the induction hypothesis. Now

$$\begin{split} s_{k+1} &= s_{\lfloor \frac{k+1}{2} \rfloor} + s_{\lfloor \frac{2k+2}{3} \rfloor} + k + 1 & \text{by definition of the sequence} \\ &> 4 \left\lfloor \frac{k+1}{2} \right\rfloor + 4 \left\lfloor \frac{2k+2}{3} \right\rfloor + k + 1 & \text{by IH and } (**) \\ &> 4 \left(\frac{k+1}{2} - 1 \right) + 4 \left(\frac{2k+2}{3} - 1 \right) + k + 1 & \text{by } (*) \\ &= 2(k-1) + \frac{4}{3}(2k-1) + k + 1 \\ &= \left(2 + \frac{4}{3} + 1 \right) k - 2 - \frac{4}{3} + 1 \\ &= \frac{17}{3}k - \frac{7}{3} \\ &= \frac{12}{3}k + 4 + \frac{5}{3}k - 4 - \frac{7}{3} \\ &= 4k + 4 + \frac{5}{3}k - \frac{19}{3} \\ &= 4(k+1) + \frac{1}{3}(5k-19) \\ &> 4(k+1) & \text{where } \frac{1}{3}(5k-19) > 0 \text{ since } k \end{split}$$

which is what we wanted to show.

Therefore, by induction, $s_n > 4n$ for all integers $n \ge 3$.

(c) Find t_2 , t_3 , t_4 , t_5 , and t_6 . Guess a formula for t_n . Solution: We see that

so we can guess that $t_n = 2^n + 1$ for all $n \ge 0$.

 $\geq 5,$

(d) Prove by strong induction that your guess in part (c) is correct for all integers $n \ge 0$.

Proof. We prove this by induction.

Base cases: (Note that first base case is n = 0.)

(n = 0): Note that $t_0 = 2$ and $2^0 + 1 = 1 + 1 = 2$. Thus $t_0 = 2^0 + 1$.

(n = 1): Note that $t_1 = 3$ and $2^1 + 1 = 2 + 1 = 3$. Thus $t_1 = 2^1 + 1$.

Inductive step: Let $k \ge 1$ be an integer. Suppose that

$$t_i = 2^i + 1$$
 for all integers $i \quad 0 \le i \le k$. (IH)

(We want to show that $t_{k+1} = 2^{k+1} + 1$.) Now

$$t_{k+1} = 3t_k - 2t_{k-1}$$
 by definition of the sequence

$$= 3 (2^k + 1) - 2 (2^{k-1} + 1)$$
 by IH

$$= 3 \cdot 2^k + 3 - 2 \cdot 2^{k-1} - 2$$

$$= 3 \cdot 2^k - 2^k + 1$$

$$= (3 - 1) \cdot 2^k + 1$$

$$= 2 \cdot 2^k + 1$$

$$= 2^{k+1} + 1,$$

which is what we wanted to show.

Therefore, by induction, $t_n = 2^n + 1$ for all integers $n \ge 0$.